

# Voting Equilibria Under Proportional Representation

## Online Appendix

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### Proof of Lemma 1

Let  $v_t = i$ . Take any  $j \in L$  and consider profile  $(j, v_{-t})$ . Since  $t$  is neither majority-pivotal nor median pivotal,  $k(j, v_{-t}) = k(v)$ , and, thus,  $p_h(j, v_{-t}) = p_h(v)$  for every  $h \in L$ . Moreover,  $k(v)$  is a majority party in  $(j, v_{-t})$  if and only if it is so in  $v$ .

Suppose  $v_t$  is not strategically sincere for  $t$  in  $v$ . There exists  $j \in L \setminus \{i\}$  such that  $u(p_j(v); t) > u(p_i(v); t)$ . Take any  $\epsilon \in (0, 1)$ . From the discussion in the previous paragraph, we conclude that

$$U(j, v_{-t}; t|\epsilon) - U(v_t, v_{-t}; t|\epsilon) \geq \frac{\epsilon}{n} [u(p_j(v); t) - u(p_i(v); t)] > 0,$$

which implies that  $v_t$  is not a robust best response to  $v_{-t}$ .

Suppose  $v_t$  is strategically sincere for  $t$  in  $v$ . Then, for every  $j \in L$ ,  $u(p_i(v); t) \geq u(p_j(v); t)$ . Thus,

$$U(v_t, v_{-t}; t|\epsilon) - U(j, v_{-t}; t|\epsilon) \geq \frac{\epsilon}{n} [u(p_i(v); t) - u(p_j(v); t)] \geq 0$$

for every  $j \in L$  and every  $\epsilon \in [0, 1)$ . Hence,  $v_t$  is a robust best response to  $v_{-t}$ . ■

## Proof of Proposition 1

1. Suppose  $|X_m| > M$ . Define the voting profile  $v$  by the following:

$$v_t = \begin{cases} m & \text{if } t \in X_m, \\ \min \left( \arg \max \left\{ u(\theta_i; t) \mid i \in \arg \max \{ u(s_j^m; t) \mid j \in L \} \right\} \right) & \text{otherwise.} \end{cases} \quad (9)$$

Since  $v_t = m$  for every  $t \in X_m$ ,  $b_m(v) > M$ . Thus,  $k(v) = m$ , and no voter is majority- or median-pivotal. By construction,  $v$  is strategically sincere. Lemma 1 then implies that  $v \in V(T, \theta, q)$ .

2. Let  $v \in V(T, \theta, q)$  and  $b_{k(v)}(v) > M$ . Since  $b_{k(v)}(v) > M$ , no voter is majority- or median-pivotal. By Lemma 1,  $v$  is strategically sincere. Suppose  $k(v) < m$ . Then  $t_M > \bar{y}_k$  which implies  $T_{k(v)}(v) \subseteq \{t_1, \dots, t_{M-1}\}$ , contradicting  $b_{k(v)}(v) > M$ . Suppose  $k(v) > m$ . Then  $t_{M-1} < \underline{y}_{k(v)}$ , which implies  $T_{k(v)}(v) \subseteq \{t_M, \dots, t_\ell\}$ , contradicting  $b_{k(v)}(v) > M$ . Thus,  $k(v) = m$ . Since  $v$  is strategically sincere,  $T_m(v) \subseteq X_m$ . Hence,  $|X_m| > M$ . ■

## Proof of Proposition 2

Suppose  $|X_m| < M - 1$  and  $\{t_{M-1}, t_M, t_{M+1}\} \subseteq X_m$ . Let  $v$  be as defined in (9). Note that, by definition of  $X_m$ ,  $m \in \arg \max \{u(s_i^m; t) \mid i \in L\}$  if and only if  $t \in X_m$ . Thus,  $T_m(v) = X_m$ , implying  $T_m(v) < M - 1$ . Also, for every  $t$  with  $v_t \neq m$ , either  $t < \underline{y}_m$  or  $t > \bar{y}_m$ . If  $t < \underline{y}_m$ , then  $t < \theta_m < s_i^m$  for every  $i > m$ . Thus, there is no  $i \in L$  such that  $i > m$  and  $i \in \arg \max \{u(s_j^m; t) \mid j \in L\}$ . Therefore,  $v_t < m$ . Similarly, if  $t > \bar{y}_m$ , then  $v_t > m$ . Then since  $\underline{y}_m \leq t_{M-1} < t_{M+1} \leq \bar{y}_m$ ,  $\sum_{i=1}^{m-1} b_i(v) < M - 1$  and  $\sum_{i=m+1}^\ell b_i(v) < M - 1$ . This implies that  $k(v) = m$  and no voter is majority- or median-pivotal. Then, by construction,  $v$  is strategically sincere, and, by Lemma 1,  $v \in V(T, \theta, q)$ . ■

### Proof of Proposition 3

1. Assume  $|X_m| = M$ . Define the voting profile  $v$  as in (9). By construction,  $T_m(v) = X_m$  and  $T^*(v) = T$ . We just have to show that  $v \in V(T, \theta, q)$ . For any  $t \in T \setminus X_m$ ,  $t$  is neither majority-pivotal, nor median-pivotal. Thus, by Lemma 1,  $v_t$  is a robust best response. Take any  $t \in X_m$  and take any  $\epsilon \in [0, 1]$ . Since  $m$  is the majority party in  $v$ ,

$$U(v; t|\epsilon) = (1 - \epsilon)u(\theta_m; t) + \frac{\epsilon}{n} \left[ Mu(\theta_m; t) + \sum_{i \in L \setminus \{m\}} b_i(v)u(s_i^m; t) \right]. \quad (10)$$

Consider voter  $t$ 's deviation by voting for some  $j \neq m$  and let  $v' = (j, v_{-t})$ . Since  $t_{M-1} \in X_m$ ,  $\sum_{i=1}^{m-1} b_i(v) < M - 1$ . Since  $t_{M+1} \in X_m$ ,  $\sum_{i=m+1}^{\ell} b_i(v) < M - 1$ . Thus,  $k(v') = m$ , implying  $p_i(v') = p_i(v) = s_i^m$  for every  $i \in L$ . Since  $m$  is not a majority party in  $v'$ ,

$$U(v'; t|\epsilon) = \frac{1}{n} \left[ (M - 1)u(\theta_m; t) + \sum_{i \in L \setminus \{m\}} b_i(v)u(s_i^m; t) + u(s_j^m; t) \right]. \quad (11)$$

Subtracting (11) from (10), we obtain

$$\begin{aligned} U(v; t|\epsilon) - U(v'; t|\epsilon) &= \frac{1 - \epsilon}{n} \left[ (M - 1)u(\theta_m; t) - \sum_{i \in L \setminus \{m\}} b_i(v)u(s_i^m; t) \right] \\ &\quad + \frac{1}{n} [u(\theta_m; t) - u(s_j^m; t)]. \end{aligned}$$

Since  $t \in X_m$ ,  $u(\theta_m; t) \geq u(s_i^m; t)$  for every  $i \in L$ . Also,  $\sum_{i \in L \setminus \{m\}} b_i(v) = M - 1$ . Hence,  $U(v; t|\epsilon) \geq U(v'; t|\epsilon)$ . Therefore,  $v$  is a robust equilibrium.

2. Assume  $|X_m| = M - 1$ . If  $m = 1$ , then  $t_1 \in X_m$ , implying  $t_{M+1} \notin X_m$ , a contradiction. If  $m = \ell$ , then  $t_n \in X_m$ , implying  $t_{M-1} \notin X_m$ , a contradiction. Thus,  $2 \leq m \leq \ell - 1$ . Let

$r_m = \frac{n-1}{2n}$ . For each  $i \in L \setminus \{m\}$ , let

$$r_i = \frac{1}{n} \left| \left\{ t \in T \setminus X_m \mid i = \min \left( \arg \max \left\{ u(\theta_i; t) \mid i \in \arg \max \{ u(s_j^m; t) \mid j \in L \} \right\} \right) \right\} \right|.$$

Note that  $\sum_{i \in L} r_i = 1$ . Either  $\sum_{i \in L} r_i s_i^m \geq \theta_m$  or  $\sum_{i \in L} r_i s_i^m < \theta_m$ . If the former is true, then let  $t^* = \max\{t \in T \mid t < \underline{y}_m\}$ . If the latter is true, then let  $t^* = \min\{t \in T \mid t > \bar{y}_m\}$ .

Note that  $t^* \notin X_m$ . Define the voting profile  $v$  by the following:

$$v_t = \begin{cases} m & \text{if } t \in X_m \cup \{t^*\}, \\ \min \left( \arg \max \left\{ u(\theta_i; t) \mid i \in \arg \max \{ u(s_j^m; t) \mid j \in L \} \right\} \right) & \text{otherwise.} \end{cases} \quad (12)$$

Note that  $T_m(v) = X_m \cup \{t^*\}$  and so  $|T_m(v)| = M$ . By construction, for every  $t \in T \setminus \{t^*\}$ ,  $v_t$  is strategically sincere in  $v$ . For any  $t \in T \setminus T_m(v)$ ,  $t$  is neither majority-pivotal, nor median-pivotal. So,  $v_t$  is a robust best response by Lemma 1. For every  $t \in X_m$ , the argument in the proof of the first statement of Proposition 3 holds true. Lastly, consider voter  $t^*$ 's deviation by voting for some  $j \neq m$ , and let  $v' = (j, v_{-t^*})$ . Since  $k(v') = m$ ,  $p_i(v') = p_i(v) = s_i^m$  for every  $i \in L$ . Let

$$i^* = \min \left( \arg \max \left\{ u(\theta_i; t^*) \mid i \in \arg \max \{ u(s_i^m; t^*) \mid i \in L \} \right\} \right).$$

Note that if  $j = i^*$ , then  $T_i(v') = nr_i$  for every  $i \in L$ ; and that if  $j \neq i^*$ , then  $T_i(v') = nr_i$  for every  $i \in L \setminus \{j, i^*\}$ ,  $T_j(v') = nr_j + 1$ , and  $T_{i^*}(v') = nr_{i^*} - 1$ . Take any  $\epsilon \in [0, 1)$ . Then,

$$\sum_{i \in L} r_i u(s_i^m; t^*) - U(v'; t^* | \epsilon) = \frac{1}{n} [u(s_{i^*}^m; t^*) - u(s_j^m; t^*)] \geq 0 \quad (13)$$

since  $u(s_{i^*}^m; t^*) = \max\{u(s_i^m; t^*) \mid i \in L\}$ . Note that, by construction, either  $t^* < \theta_m \leq \sum_{i \in L} r_i s_i^m$  or  $\sum_{i \in L} r_i s_i^m < \theta_m < t^*$ . Then, since  $f$  is strictly concave,  $u(\theta_m; t^*) > \sum_{i \in L} r_i u(s_i^m; t^*)$ .

Then, from (13), we conclude that  $u(\theta_m; t^*) > U(v'; t^*|\epsilon)$ . Then, for sufficiently small  $\epsilon$ ,

$$U(v; t^*|\epsilon) - U(v'; t^*|\epsilon) = (1-\epsilon)[u(\theta_m; t^*) - U(v'; t^*|\epsilon)] + \epsilon \left[ \sum_{i \in L} \frac{b_i(v)}{n} u(s_i^m; t^*) - U(v'; t^*|\epsilon) \right] > 0.$$

Thus,  $v$  is a robust equilibrium of  $G(T, \theta, q, 0)$ . ■

## Proof of Proposition 4

Assume  $\{t_{M-1}, t_M, t_{M+1}\} \not\subseteq X_m$ . For each  $t \in T$ , let

$$\alpha(t) = \min \left( \arg \max \left\{ u(\theta_i; t) \mid i \in \arg \max \left\{ u(s_j^{m+1}; t) \mid j \in \arg \max \{u(s_h^m; t) \mid h \in L\} \right\} \right\} \right),$$

and define voting profile  $\hat{v}$  by the following.

$$\hat{v}_t = \begin{cases} m & \text{if } t \in [\underline{y}_m, t_M], \\ \alpha(t) & \text{otherwise;} \end{cases} \quad (14)$$

Recall that  $\theta_1 < t_M < \theta_\ell$  and  $\theta_m \leq t_M$ . This implies that  $m \leq \ell - 1$ . For every  $t > \max\{t_M, \bar{y}_m\}$ ,  $\hat{v}_t = \alpha(t) \geq m + 1$ . If  $m = 1$ , then  $\underline{y}_m = t_1$ , so  $\{t_1, \dots, t_M\} \subseteq T_m(\hat{v})$ . Otherwise, for every  $t < \underline{y}_m$ ,  $\hat{v}_t = \alpha(t) \leq m - 1$ . Thus, when  $m = 1$ , party  $m$  is the majority party, and when  $m > 1$ , party  $m$  is the median party. Then, for every  $i \in L$ ,  $p_i(\hat{v}) = s_i^m$ . For every  $\epsilon \in [0, 1]$  and every  $t \in T$ ,

$$U(\hat{v}; t|\epsilon) = (1 - \epsilon)u(\theta_1; t) + \frac{\epsilon}{n} \sum_{i \in L} b_i(\hat{v})u(s_i^m; t) \quad (15)$$

if  $m = 1$ ; and

$$U(\hat{v}; t|\epsilon) = \frac{1}{n} \sum_{i \in L} b_i(\hat{v})u(s_i^m; t) \quad (16)$$

if  $m > 1$ .

The proof will include a series of lemmas. The first lemma shows that, for voters who vote for the median party or any party to the right of the median in profile  $\hat{v}$ , a deviation by voting for any party to the left of the median is not profitable.

**Lemma 3** *Assume  $m > 1$ . For every  $t \geq \underline{y}_m$  and every  $j \leq m - 1$ , there exists  $\bar{\epsilon} > 0$  such that  $U(\hat{v}; t|\epsilon) \geq U(j, \hat{v}_{-t}; t|\epsilon)$  for every  $\epsilon \in [0, \bar{\epsilon}]$ .*

*Proof:* Take any  $t \geq \underline{y}_m$  and let  $h = \hat{v}_t$ . Note that  $h \geq m$ . Take any  $j \leq m - 1$ . First, suppose  $t_{M-1} \geq \underline{y}_m$ . Then  $\sum_{i=1}^{m-1} b_i(\hat{v}) < M - 1$ , implying  $k(j, \hat{v}_{-t}) = m$ . Then, for every  $\epsilon \in [0, 1]$ ,

$$U(\hat{v}; t|\epsilon) - U(j, \hat{v}_{-t}; t|\epsilon) = \frac{1}{n}[u(s_h^m; t) - u(s_j^m; t)]. \quad (17)$$

If  $t \geq t_{M+1}$ , then  $h = \alpha(t) \in \arg \max\{u(s_i^m; t) | i \in L\}$ . So,  $u(s_h^m; t) \geq u(s_j^m; t)$ , implying (17) is nonnegative. If  $t \in [\underline{y}_m, t_M]$ , then  $h = m$ . Since  $t \geq \underline{y}_m = \frac{s_{m-1}^m + s_m^m}{2}$ ,  $u(s_m^m; t) \geq u(s_j^m; t)$ . Thus, (17) is nonnegative.

Now suppose  $t_{M-1} < \underline{y}_m$ . Then,  $\sum_{i=1}^{m-1} b_i(\hat{v}) = M - 1$ . A1 implies that, for each  $i = 2, \dots, \ell$ ,  $(\frac{\theta_{i-1} + \theta_i}{2}, \theta_i) \cap T \neq \emptyset$ . Since  $t_M \geq \theta_m$ , it must be that  $t_{M-1} > \frac{\theta_{m-1} + \theta_m}{2}$ . Since  $\underline{y}_m > t_{M-1}$ ,  $\underline{y}_m > \frac{\theta_{m-1} + \theta_m}{2}$ , which implies  $\underline{x}_m(q) > \theta_{m-1}$ . Then, it must be that  $q > \theta_{m-1}$ . Since  $t_{M-1} \in (\frac{\theta_{m-1} + \theta_m}{2}, \underline{y}_m)$ ,  $\alpha(t_{M-1}) = m - 1$ , so  $b_{m-1}(\hat{v}) > 0$ . This implies that  $k(j, \hat{v}_{-t}) = m - 1$ , and, so,  $p_i(j, \hat{v}_{-t}) = s_i^{m-1}$  for every  $i \in L$ . We consider two cases separately:  $m > 2$  and  $m = 2$ .

First, suppose  $m > 2$ . Since  $t_1 < \theta_1$ ,  $\alpha(t_1) = 1$ . So,  $b_1(\hat{v}) > 0$ , implying  $b_{m-1}(j, \hat{v}_{-t}) < M$ . Then, for every  $\epsilon \in [0, 1]$ ,

$$U(j, \hat{v}_{-t}; t|\epsilon) = \frac{1}{n} \left( \sum_{i \in L} b_i(\hat{v}) u(s_i^{m-1}; t) + [u(s_j^{m-1}; t) - u(s_h^{m-1}; t)] \right). \quad (18)$$

Suppose  $q > \theta_m$ , i.e.,  $\underline{x}_m(q) = 2\theta_m - q$ . Then,  $A(j, \hat{v}_{-t}) = [2\theta_{m-1} - q, q]$ . Since  $\underline{x}_m(q) \in (\theta_{m-1}, \theta_m)$ , for every  $i \geq m$ ,  $s_i^m = s_i^{m-1}$ . For every  $i \leq m - 1$ ,  $s_i^m = 2\theta_m - q > s_i^{m-1}$ . Since  $t \geq t_M > 2\theta_m - q$ ,  $u(2\theta_m - q; t) > u(s_i^{m-1}; t)$  for every  $i \leq m - 1$ . Also, since  $t > \theta_m$ ,

$s_j^{m-1} < 2\theta_m - q$ , and  $s_h^{m-1} \in [\theta_m, q]$ , it must be the case that  $u(s_h^{m-1}; t) > u(s_j^{m-1}; t)$ . Then,

$$\begin{aligned} U(\hat{v}; t|\epsilon) - U(j, \hat{v}_{-t}; t|\epsilon) &= \frac{1}{n} \left( \sum_{i=1}^{m-1} b_i(\hat{v}) [u(2\theta_m - q; t) - u(s_i^{m-1}; t)] + [u(s_h^{m-1}; t) - u(s_j^{m-1}; t)] \right) \\ &> 0 \end{aligned}$$

for every  $\epsilon \in [0, 1]$ . Suppose  $q \in (\theta_{m-1}, \theta_m)$ . Then, for every  $i \geq m$ ,  $s_i^{m-1} = q$ , and, for every  $i \leq m-1$ ,  $s_i^m = q$ . Thus, for every  $\epsilon \in [0, 1]$ ,

$$\begin{aligned} U(\hat{v}; t|\epsilon) - U(j, \hat{v}_{-t}; t|\epsilon) &= \frac{1}{n} \left( \sum_{i=1}^{m-1} b_i(\hat{v}) [u(q; t) - u(s_i^{m-1}; t)] \right. \\ &\quad \left. + \sum_{i=m}^{\ell} b_i(\hat{v}) [u(s_i^m; t) - u(q; t)] + [u(q; t) - u(s_j^{m-1}; t)] \right) \end{aligned} \quad (19)$$

Since  $s_i^{m-1} < q < t$ ,  $u(q; t) > u(s_i^{m-1}; t)$  for every  $i \leq m-1$  (including  $j$ ). Since  $t > \theta_m$  and  $s_i^m \in [\theta_m, 2\theta_m - q]$ ,  $u(s_i^m; t) > u(q; t)$  for every  $i \geq m$ . Thus, (19) is positive.

Now suppose  $m = 2$ . Then,  $b_1(\hat{v}) = M - 1$  and  $j = 1$ . Thus, party 1 is the majority party in  $(j, \hat{v}_{-t})$ . Let  $C(t) = \frac{1}{n} \sum_{i \in L} b_i(\hat{v}) u(s_i^m; t) - u(\theta_1; t)$  and let

$$G(t) = \frac{1}{n} \left( \sum_{i \in L} b_i(\hat{v}) u(s_i^m; t) - \sum_{i \in L} b_i(j, \hat{v}_{-t}) u(s_i^{m-1}; t) \right).$$

Then, for every  $\epsilon \in [0, 1]$ ,

$$U(\hat{v}; t|\epsilon) - U(j, \hat{v}_{-t}; t|\epsilon) = (1 - \epsilon)C(t) + \epsilon G(t). \quad (20)$$

Note that  $t > \theta_m$ ,  $s_i^m \in [\underline{x}_m(q), \bar{x}_m(q)]$ , and  $\theta_1 < \underline{x}_m(q)$ . This implies that  $u(s_i^m; t) > u(\theta_1; t)$  for every  $i \in L$ . Hence,  $C(t) > 0$ . If  $G(t) \geq 0$ , then (20) is positive for every  $\epsilon \in [0, 1]$ . If  $G(t) < 0$ , then let  $\bar{\epsilon} = \frac{C(t)}{C(t) - G(t)}$ . Then, for every  $\epsilon \in [0, \bar{\epsilon}]$ ,  $U(\hat{v}; t|\epsilon) \geq U(j, \hat{v}_{-t}; t|\epsilon)$ , which completes the proof of the lemma.  $\blacksquare$

We are ready to prove the first statement of the proposition. Suppose  $t_{M+1} \in X_m$ . Since  $\theta_m \leq t_M < t_{M+1} \leq \bar{y}_m$ ,  $t_M \in X_m$ . Then, it must be the case that  $t_{M-1} < \underline{y}_m$ , implying  $m > 1$ . Note that  $[\underline{y}_m, t_M] \cap T = \{t_M\}$ . Since  $t_M \in X_m$ ,  $\hat{v}_{t_M} = m \in \arg \max\{u(s_i^m; t_M) | i \in L\}$ . Also, by construction,  $\hat{v}_t = \alpha(t) \in \arg \max\{u(s_i^m; t) | i \in L\}$  for every  $t \in T \setminus \{t_M\}$ . Thus,  $\hat{v}$  is strategically sincere, i.e.,  $T^*(\hat{v}) = T$ .

I now will show that  $\hat{v} \in V(T, \theta, q)$ . Since  $t_{M+1} \in X_m$ ,  $\hat{v}_{t_{M+1}} = \alpha(t_{M+1}) = m$ . This implies that  $\sum_{i=m+1}^{\ell} b_i(\hat{v}) < M - 1$ . Then, for every  $t \leq t_{M-1}$ , voter  $t$  is neither majority-pivotal nor median-pivotal. Thus, by Lemma 1,  $\hat{v}_t$  is a robust best response to  $\hat{v}_{-t}$  for every  $t \leq t_{M-1}$ . Take any  $t \geq t_M$ , and consider voter  $t$ 's deviation by voting for any  $j \neq \hat{v}_t$ . If  $j \geq t_M$ , the deviation would not change the median party, i.e.,  $k(j, \hat{v}_{-t}) = m$ . Then, for every  $i \in L$ ,  $p_i(j, \hat{v}_{-t}) = s_i^m = p_i(\hat{v})$ . Since  $\hat{v}_t \in \arg \max\{u(s_i^m; t) | i \in L\}$ ,  $U(\hat{v}; t|\epsilon) \geq U(j, \hat{v}_{-t}; t|\epsilon)$  for every  $\epsilon \in [0, 1]$ . Suppose  $j \leq m - 1$ . By Lemma 3, there exists  $\bar{\epsilon} > 0$  such that  $U(\hat{v}; t|\epsilon) \geq U(j, \hat{v}_{-t}; t|\epsilon)$  for every  $\epsilon \in [0, \bar{\epsilon}]$ . Thus,  $\hat{v} \in V(T, \theta, q)$ , which completes the proof of the first statement in Proposition 4.

To prove the second statement, assume  $t_{M+1} \notin X_m$ . Let

$$\beta(t) = \min \left( \arg \max \left\{ u(\theta_i; t) \mid i \in \arg \max \left\{ u(s_j^m; t) \mid j \in \arg \max \{ u(s_h^{m+1}; t) \mid h \in L \} \right\} \right\} \right).$$

Define voting profile  $\tilde{v}$  by

$$\tilde{v}_t = \begin{cases} m + 1 & \text{if } t \in [t_M, \bar{y}_{m+1}], \\ \beta(t) & \text{otherwise.} \end{cases} \quad (21)$$

I will prove that either  $\hat{v}$  or  $\tilde{v}$  is a robust voting equilibrium of  $G(T, \theta, q, 0)$ . Define voting profiles  $\hat{v}'$  and  $\tilde{v}'$  by the following.

$$\hat{v}'_t = \begin{cases} m + 1 & \text{if } t = t_M, \\ \hat{v}_t & \text{otherwise;} \end{cases} \quad (22)$$



and

$$\tilde{v}'_t = \begin{cases} m & \text{if } t = t_M, \\ \tilde{v}_t & \text{otherwise.} \end{cases} \quad (23)$$

That is,  $\hat{v}'$  is the voting profile in which the median voter unilaterally deviates from  $\hat{v}$  by voting for  $m + 1$ , and  $\tilde{v}'$  is the voting profile in which the median voter unilaterally deviates from  $\tilde{v}$  by voting for  $m$ . For each  $t \in T$  and each  $\epsilon \in [0, 1]$ , let  $\hat{\Delta}(t|\epsilon) = U(\hat{v}; t|\epsilon) - U(\hat{v}'; t|\epsilon)$  and  $\tilde{\Delta}(t|\epsilon) = U(\tilde{v}; t|\epsilon) - U(\tilde{v}'; t|\epsilon)$ .

Note that, for every  $t \leq \underline{y}_m$ ,  $\alpha(t) \leq m$ . Also, since  $t_{M+1} \geq \bar{y}_m$ , for every  $t \geq t_{M+1}$ ,  $\alpha(t) \geq m + 1$ . Thus,

$$\sum_{i=1}^m b_i(\hat{v}) = M \quad \text{and} \quad \sum_{i=1}^m b_i(\hat{v}') = M - 1. \quad (24)$$

Since  $t_{M-1} < \frac{\theta_m + \theta_{m+1}}{2} \leq \frac{s_m^{m+1} + s_{m+1}^{m+1}}{2}$ ,  $\beta(t) \leq m$  for every  $t \leq t_{M-1}$ . Clearly, for every  $t \geq \bar{y}_{m+1}$ ,  $\beta(t) \geq m + 1$ . Hence,

$$\sum_{i=1}^m b_i(\tilde{v}) = M - 1 \quad \text{and} \quad \sum_{i=1}^m b_i(\tilde{v}') = M. \quad (25)$$

An implication of (24) and (25) is that  $k(\hat{v}) = k(\tilde{v}') = m$  and  $k(\hat{v}') = k(\tilde{v}) = m + 1$ . Thus, for every  $i \in L$ ,  $p_i(\hat{v}) = p_i(\tilde{v}') = s_i^m$  and  $p_i(\hat{v}') = p_i(\tilde{v}) = s_i^{m+1}$ . I now present a series of lemmas.

**Lemma 4** *For each given  $\epsilon \in [0, 1]$ ,  $\hat{\Delta}(t|\epsilon)$  is decreasing in  $t$  and  $\tilde{\Delta}(t|\epsilon)$  is increasing in  $t$ .*

*Proof:* For each  $t \in T$ , let

$$\hat{D}(t) = \frac{1}{n} \left[ \sum_{i \in L} b_i(\hat{v}) u(s_i^m; t) - \sum_{i \in L} b_i(\hat{v}') u(s_i^{m+1}; t) \right] \quad (26)$$

and

$$\tilde{D}(t) = \frac{1}{n} \left[ \sum_{i \in L} b_i(\tilde{v}) u(s_i^{m+1}; t) - \sum_{i \in L} b_i(\tilde{v}') u(s_i^m; t) \right] \quad (27)$$

I first claim  $\hat{D}$  is decreasing and  $\tilde{D}$  is increasing in  $t$ . Since  $\hat{v}$  and  $\hat{v}'$  differ only in that  $\hat{v}_{t_M} = m$  and  $\hat{v}'_{t_M} = m + 1$ , we write

$$\begin{aligned} \hat{D}(t) &= \frac{1}{n} \left( \sum_{i \in L} b_i(\hat{v}) [u(s_i^m; t) - u(s_i^{m+1}; t)] + [u(s_m^{m+1}; t) - u(s_{m+1}^{m+1}; t)] \right) \\ &= \frac{1}{n} \left( \sum_{i \in L} b_i(\hat{v}) [f(|s_i^m - t|) - f(|s_i^{m+1} - t|)] + [f(|s_m^{m+1} - t|) - f(|s_{m+1}^{m+1} - t|)] \right) \end{aligned} \quad (28)$$

Note that, for each  $i \in L$ ,  $s_i^m \leq s_i^{m+1}$  and  $s_m^{m+1} \leq s_{m+1}^{m+1}$ . Then, since  $f$  is decreasing and concave, for each  $i \in L$ ,  $f(|s_i^m - t|) - f(|s_i^{m+1} - t|)$  is decreasing in  $t$  and  $f(|s_m^{m+1} - t|) - f(|s_{m+1}^{m+1} - t|)$  is decreasing in  $t$ . Hence  $\hat{D}$  is decreasing in  $t$ . A symmetric argument proves that  $\tilde{D}$  is increasing in  $t$ .

Let  $\epsilon \in [0, 1]$ . First, suppose that  $b_m(\hat{v}) < M$  and  $b_{m+1}(\hat{v}') < M$ . Then,  $\hat{\Delta}(t|\epsilon) = \hat{D}(t)$ , implying  $\hat{\Delta}(t|\epsilon)$  is decreasing in  $t$ . Second, suppose  $b_m(\hat{v}) = M$ . Since  $t_1 < \theta_1$ ,  $\hat{v}_{t_1} = 1$ . This, together with (24), implies that  $m = 1$ . Since  $t_n \geq \theta_\ell$ ,  $\hat{v}_{t_n} = \ell$ , implying  $b_{m+1}(\hat{v}') < M$ . Then,

$$\hat{\Delta}(t|\epsilon) = (1 - \epsilon) \left[ u(\theta_1; t) - \frac{1}{n} \sum_{i \in L} b_i(\hat{v}') u(s_i^{m+1}; t) \right] + \epsilon \hat{D}(t).$$

But since  $\theta_1 \leq s_i^{m+1}$  for every  $i \in L$ , the expression in the square bracket is decreasing in  $t$ . Thus,  $\hat{\Delta}(t|\epsilon)$  is decreasing in  $t$ . Lastly, suppose  $b_{m+1}(\hat{v}') = M$ . Again since  $\hat{v}_{t_n} = \hat{v}'_{t_n} = \ell$ , it must be the case that  $m + 1 = \ell$ . Then since  $\hat{v}_{t_1} = 1$ ,  $b_m(\hat{v}) < M - 1$ . Then,

$$\hat{\Delta}(t|\epsilon) = (1 - \epsilon) \left[ \frac{1}{n} \sum_{i \in L} b_i(\hat{v}) u(s_i^m; t) - u(\theta_\ell; t) \right] + \epsilon \hat{D}(t).$$

But since  $s_i^m \leq \theta_\ell$  for every  $i \in L$ , the expression in the square bracket is decreasing in  $t$ . Thus,  $\hat{\Delta}(t|\epsilon)$  is decreasing in  $t$ . A symmetric argument proves  $\tilde{\Delta}(t|\epsilon)$  is increasing in  $t$ . ■

**Lemma 5** *The following is true.*

1. If  $\hat{\Delta}(t_M|0) > 0$ , then  $\hat{v}$  is a robust equilibrium of  $G(T, \theta, q, 0)$ .
2. If  $\tilde{\Delta}(t_M|0) > 0$ , then  $\tilde{v}$  is a robust equilibrium of  $G(T, \theta, q, 0)$ .
3. If  $\hat{\Delta}(t_M|0) = \tilde{\Delta}(t_M|0) = 0$ , then either  $\hat{v}$  or  $\tilde{v}$  is a robust equilibrium of  $G(T, \theta, q, 0)$ .

*Proof:* 1. Suppose  $\hat{\Delta}(t_M) > 0$ . Take any  $t \in T$ , and let  $h = \hat{v}_t$ . Assume  $t \geq t_{M+1}$  and notice that  $\hat{v}_t \geq m + 1$ . Consider voter  $t$ 's deviation by voting for  $j$ . Suppose  $j \geq m$ . Then the deviation does not change the majority or the median status of party  $m$ . Since  $h \in \arg \max\{u(s_i^m; t) | i \in L\}$ ,  $U(\hat{v}; t|\epsilon) \geq U(j, \hat{v}_{-t}; t|\epsilon)$  for every  $\epsilon \in [0, 1]$ . Suppose  $j \leq m - 1$ . Then, by Lemma 3,  $U(\hat{v}; t|\epsilon) \geq U(j, \hat{v}_{-t}; t|\epsilon)$  for sufficiently small  $\epsilon$ . Thus,  $\hat{v}_t$  is a robust best response to  $\hat{v}_{-t}$ .

Assume  $t \leq t_M$ . Again, consider voter  $t$ 's deviation from  $\hat{v}$  by voting for any  $j \neq h$ , i.e., we consider the profile  $(j, \hat{v}_{-t})$ . If  $j \leq m$ , then the deviation would not change the identity of the median or majority party. So,  $k(j, \hat{v}_{-t}) = k(\hat{v}) = m$ , and, for every  $i \in L$ ,  $p_i(j, \hat{v}_{-t}) = p_i(\hat{v}) = s_i^m$ . And since, by construction,  $h \in \arg \max\{u(s_i^m; t) | i \leq m\}$ ,  $U(j, \hat{v}_{-t}; t|\epsilon) \leq U(\hat{v}; t|\epsilon)$  for every  $\epsilon \in [0, 1]$ .

Now suppose  $j \geq m + 1$ . Then  $k(j, \hat{v}_{-t}) = m + 1$  and  $p_i(j, \hat{v}_{-t}) = s_i^{m+1}$  for every  $i \in L$ . Note that the only possible difference between  $\hat{v}'$  and  $(j, \hat{v}_{-t})$  is that, in  $(j, \hat{v}_{-t})$ , one vote for  $h$  in  $\hat{v}$  is transferred to  $j$ , and, in  $\hat{v}'$ , one vote for  $m$  is transferred to  $m + 1$ . I claim  $U(\hat{v}'; t|\epsilon) \geq U(j, \hat{v}_{-t}; t|\epsilon)$  for every  $\epsilon \in [0, 1]$ . To see this, first, suppose  $m < \ell - 1$ . Then, there is no majority party in  $\hat{v}'$  or  $(j, \hat{v}_{-t})$ . So, for every  $\epsilon \in [0, 1]$ ,

$$U(\hat{v}'; t|\epsilon) - U(j, \hat{v}_{-t}; t|\epsilon) = \frac{1}{n} \left( [u(s_{m+1}^{m+1}; t) - u(s_j^{m+1}; t)] + [u(s_h^{m+1}; t) - u(s_m^{m+1}; t)] \right).$$

Since  $t \leq t_M$  and  $m + 1 \leq j$ , we have  $t < \theta_{m+1} = s_{m+1}^{m+1} \leq s_j^{m+1}$ , implying  $u(s_{m+1}^{m+1}; t) \geq u(s_j^{m+1}; t)$ . If  $h = m$ , then clearly  $u(s_h^{m+1}; t) = u(s_m^{m+1}; t)$ . Suppose  $h < m$ . Then, since

$h = \alpha(t)$ ,  $t \leq \frac{s_h^m + \theta_m}{2}$ . But since  $s_h^m \leq s_h^{m+1} \leq s_m^{m+1}$  and  $\theta_m \leq s_m^{m+1}$ ,  $\frac{s_h^m + \theta_m}{2} \leq \frac{s_h^{m+1} + s_m^{m+1}}{2}$ , implying  $u(s_h^{m+1}; t) \geq u(s_m^{m+1}; t)$ . Therefore,  $U(\hat{v}'; t|\epsilon) \geq U(j, \hat{v}_{-t}; t|\epsilon)$ . Second, suppose  $m = \ell - 1$ . Then,  $m + 1$  is the majority party in  $\hat{v}'$  and  $(j, \hat{v}_{-t})$ , and  $j = m + 1 = \ell$ . Then,

$$U(\hat{v}'; t|\epsilon) - U(j, \hat{v}_{-t}; t|\epsilon) = \frac{\epsilon}{n}[u(s_h^{m+1}; t) - u(s_m^{m+1}; t)] \geq 0.$$

Hence, the claim is true. This implies that, if  $U(\hat{v}; t|\epsilon) \geq U(\hat{v}'; t|\epsilon)$  for sufficiently small  $\epsilon$ , then  $\hat{v}_t$  is a robust best response to  $\hat{v}_{-t}$ . Thus, it suffices to show that  $\hat{\Delta}(t|\epsilon) \geq 0$  for sufficiently small  $\epsilon$ . But since  $t \leq t_M$  and  $\hat{\Delta}(t|\epsilon)$  is decreasing in  $t$  by Lemma 4, it suffices to show  $\hat{\Delta}(t_M|\epsilon) \geq 0$  for sufficiently small  $\epsilon$ . Suppose  $1 < m < \ell - 1$ . Then, for every  $\epsilon \in [0, 1]$ ,  $\hat{\Delta}(t_M|\epsilon) = \hat{\Delta}(t_M|0) > 0$ . Suppose  $m = 1$  or  $m = \ell - 1$ . Then,

$$\hat{\Delta}(t_M|\epsilon) = (1 - \epsilon)\hat{\Delta}(t_M|0) + \epsilon\hat{D}(t_M). \quad (29)$$

If  $\hat{D}(t_M)$  is nonnegative, then  $\hat{\Delta}(t_M|\epsilon) \geq 0$  for every  $\epsilon \in [0, 1]$ . If  $\hat{D}(t_M) < 0$ , then  $\hat{\Delta}(t_M|\epsilon) \geq 0$  for every  $\epsilon \in [0, \frac{\hat{\Delta}(t_M|0)}{\hat{\Delta}(t_M|0) - \hat{D}(t_M)}]$ . Therefore,  $\hat{v} \in V(T, \theta, q)$ .

2. A symmetric argument proves the second statement.

3. Suppose  $\hat{\Delta}(t_M|0) = \tilde{\Delta}(t_M|0) = 0$ . Again, note that, if  $\hat{\Delta}(t_M|\epsilon) \geq 0$  for sufficiently small  $\epsilon$ , then  $\hat{v}$  is a robust equilibrium, and that, if  $\tilde{\Delta}(t_M|\epsilon) \geq 0$  for sufficiently small  $\epsilon$ , then  $\tilde{v}$  is a robust equilibrium. If  $1 < m < \ell - 1$ , then  $\hat{\Delta}(t_M|\epsilon) = \hat{\Delta}(t_M|0) = 0$  for every  $\epsilon \in [0, 1]$ . Thus,  $\hat{v}$  is a robust equilibrium. Suppose  $m = 1$ . Since  $\hat{\Delta}(t_M|0) = 0$ , we obtain from (29) that  $\hat{\Delta}(t_M|\epsilon) = \epsilon\hat{D}(t_M)$ . Similarly, because  $\tilde{\Delta}(t_M|0) = 0$ ,  $\tilde{\Delta}(t_M|\epsilon) = \epsilon\tilde{D}(t_M)$ . So, it suffices to prove that either  $\hat{D}(t_M) \geq 0$  or  $\tilde{D}(t_M) \geq 0$ .

Note that

$$\hat{\Delta}(t_M|0) = u(\theta_1; t_M) - \frac{1}{n} \sum_{i \in L} b_i(\hat{v}')u(s_i^2; t_M) \quad (30)$$

and

$$\tilde{\Delta}(t_M|0) = \frac{1}{n} \sum_{i \in L} b_i(\tilde{v})u(s_i^2; t_M) - u(\theta_1; t_M). \quad (31)$$

Since  $\hat{\Delta}(t_M|0) = \tilde{\Delta}(t_M|0) = 0$ , we have

$$\frac{1}{n} \sum_{i \in L} b_i(\hat{v}')u(s_i^2; t_M) = \frac{1}{n} \sum_{i \in L} b_i(\tilde{v})u(s_i^2; t_M) = u(\theta_1; t_M).$$

Then, from (26) and (27), we obtain that

$$\hat{D}(t_M) = \frac{1}{n} \sum_{i \in L} b_i(\hat{v})u(s_i^1; t_M) - u(\theta_1; t_M) \quad (32)$$

and

$$\tilde{D}(t_M) = u(\theta_1; t_M) - \frac{1}{n} \sum_{i \in L} b_i(\tilde{v}')u(s_i^1; t_M). \quad (33)$$

Since  $t_M \leq \frac{\theta_1 + \theta_2}{2}$ ,  $u(\theta_1; t_M) \geq u(s_2^2; t_M)$  and, for every  $i \geq 3$ ,  $u(\theta_1; t_M) > u(s_i^2; t_M)$ . Then since  $\hat{\Delta}(t_M|0) = 0$ , (30) implies  $u(\theta_1; t_M) < u(s_1^2; t_M)$ . Then, it must be that  $s_1^2 = \underline{x}_2(q) \in (\theta_1, \theta_2)$ . Suppose  $\underline{x}_2(q) = q$ . Then, for every  $i \geq 2$ ,  $s_i^1 = q$ . And since  $u(q; t_M) > u(\theta_1; t_M)$ , we conclude  $\hat{D}(t_M) > 0$  from (32). This implies  $\hat{\Delta}(t_M|\epsilon) \geq 0$  for every  $\epsilon \in [0, 1]$ . Hence,  $\hat{v} \in V(T, \theta, q)$ . Now suppose  $\underline{x}_2(q) = 2\theta_2 - q$ . This implies  $q > \theta_2$ . Comparing the definitions of  $\hat{v}$  and  $\tilde{v}'$ , we first conclude that  $T_1(\hat{v}) = T_1(\tilde{v}') = \{t_1, \dots, t_M\}$ . Note that, since  $q > \theta_2$ ,  $\bar{x}_1(q) = \bar{x}_2(q) = q$ , which implied that, for every  $i \geq 2$ ,  $s_i^1 = s_i^2$ . Consider a voter  $t \in [t_{M+1}, \bar{y}_2]$ . By the definition of  $\tilde{v}'$ ,  $\tilde{v}'_t = 2$ . Since  $t \geq t_{M+1} > \bar{y}_1 = \frac{\theta_1 + \theta_2}{2}$ ,  $u(s_1^1; t) = u(\theta_1; t) < u(\theta_2; t) = u(s_2^1; t)$ . Since  $t \leq \bar{y}_2$ ,  $u(s_2^1; t) = u(s_2^2; t) \geq u(s_3^2; t) = u(s_3^1; t)$ . Thus,  $\hat{v}_t = \alpha(t) = 2$  as well. Since  $s_i^1 = s_i^2$  for every  $i \geq 3$ ,  $\alpha(t) = \beta(t)$  for every  $t > \bar{y}_2$ . Thus, we conclude that  $\hat{v}_t = \tilde{v}'_t$  for every  $t \in T$ , which implies  $b_i(\hat{v}) = b_i(\tilde{v}')$  for every  $i \in L$ . Then, from (32) and (33), we conclude that either  $\hat{D}(t_M) \geq 0$  or  $\tilde{D}(t_M) \geq 0$ . Thus, either  $\hat{v}$  or  $\tilde{v}$  is a robust equilibrium. A symmetric argument proves the statement for the case that  $m = \ell - 1$ . ■

**Lemma 6**  $\hat{\Delta}(t_M|0) + \tilde{\Delta}(t_M|0) \geq 0$ .

*Proof:* We consider three mutually exclusive and jointly exhaustive cases.

CASE 1: Assume  $m = 1$ .

Note that party 1 is the majority party in  $\hat{v}$  and  $\tilde{v}'$ , and party 2 is the median party in  $\hat{v}'$  and  $\tilde{v}$ . Then, by definition,

$$\hat{\Delta}(t_M|0) = u(\theta_1; t_M) - \frac{1}{n} \sum_{i \in L} b_i(\hat{v}') u(s_i^2; t_M) \quad (34)$$

and

$$\tilde{\Delta}(t_M|0) = \frac{1}{n} \sum_{i \in L} b_i(\tilde{v}) u(s_i^2; t_M) - u(\theta_1; t_M). \quad (35)$$

By adding (34) and (35), we write

$$\hat{\Delta}(t_M|0) + \tilde{\Delta}(t_M|0) = \frac{1}{n} \sum_{i \in L} [b_i(\tilde{v}) - b_i(\hat{v}')] u(s_i^2; t_M). \quad (36)$$

From (21), we write

$$\tilde{v}_t = \begin{cases} 2 & \text{if } t \in [t_M, \bar{y}_2], \\ \beta(t) & \text{otherwise.} \end{cases} \quad (37)$$

Also, from (14) and (22), we write

$$\hat{v}'_t = \begin{cases} 1 & \text{if } t \leq t_{M-1} \\ 2 & \text{if } t = t_M \\ \alpha(t) & \text{otherwise.} \end{cases} \quad (38)$$

Since  $t_{M-1} < \frac{\theta_1 + \theta_2}{2}$  and  $\theta_1 \leq s_1^2 < s_2^2 = \theta_2$ , for every  $t \leq t_{M-1}$ ,  $\arg \max\{u(s_i^2; t) | i \in L\} = \{1\}$ . Thus, for every  $t \leq t_{M-1}$ ,  $\tilde{v}_t = \beta(t) = 1$ . For any  $t > \bar{y}_2$ , clearly  $\beta(t) \neq 1$ . Thus,  $T_1(\tilde{v}) = \{t_1, \dots, t_{M-1}\}$ . Also, since  $t_{M+1} > \bar{y}_1$ , for any  $t \geq t_{M+1}$ ,  $\alpha(t) \neq 1$ . So,

$T_1(\hat{v}') = \{t_1, \dots, t_{M-1}\}$ . Therefore,  $b_1(\tilde{v}) = b_1(\hat{v}')$ . Then, (36) is reduced to

$$\Delta(t_M|0) + \tilde{\Delta}(t_M|0) = \frac{1}{n} \sum_{i=2}^{\ell} [b_i(\tilde{v}) - b_i(\hat{v}')] u(s_i^2; t_M). \quad (39)$$

First, suppose  $q < \theta_1$ . Let  $L^- = \{i \in L | \theta_i < 2\theta_1 - q\}$ . Suppose  $L^- = L$ . Then,  $s_i^1 = s_i^2 = \theta_i$  for every  $i \in L$ . Then,  $\alpha(t) = 2$  if and only if  $t \in (\frac{\theta_1 + \theta_2}{2}, \frac{\theta_2 + \theta_3}{2}] = (\bar{y}_1, \bar{y}_2]$ . Since  $t_{M+1} > \bar{y}_1$ ,  $T_2(\tilde{v}) = T_2(\hat{v}')$ . Also, since  $s_i^1 = s_i^2 = \theta_i$  for every  $i \in L$ ,  $\alpha(t) = \beta(t)$  for every  $t \in L$ . Hence,  $b_i(\tilde{v}) = b_i(\hat{v}')$  for every  $i = 2, \dots, \ell$ . Therefore,  $\hat{\Delta}(t_M|0) + \tilde{\Delta}(t_M|0) = 0$ . Suppose  $L^- \neq L$ . Let  $\underline{j} = \max L^-$ .

Suppose  $\underline{j} \geq 2$ . For every  $t \leq \frac{\theta_{\underline{j}} + 2\theta_1 - q}{2}$ ,

$$\alpha(t) = \beta(t) = \min \left( \arg \max \{u(\theta_i; t) | i = 1, \dots, \underline{j}\} \right).$$

For every  $t > \frac{\theta_{\underline{j}} + s_{\underline{j}+1}^2}{2}$ ,

$$\alpha(t) = \beta(t) = \min \left( \arg \max \left\{ u(\theta_i; t) | i \in \arg \max \{u(s_j^2; t) | j = \underline{j} + 1, \dots, \ell\} \right\} \right).$$

Let  $\tilde{T} = \{t \in T | \frac{\theta_{\underline{j}} + 2\theta_1 - q}{2} < t \leq \frac{\theta_{\underline{j}} + s_{\underline{j}+1}^2}{2}\}$ . For every  $t \in \tilde{T}$ ,  $\hat{v}'_t = \underline{j} + 1$  and  $\tilde{v}_t = \underline{j}$ . Hence,

$$\hat{\Delta}(t_M|0) + \tilde{\Delta}(t_M|0) = \frac{|\tilde{T}|}{n} [u(\theta_{\underline{j}}; t_M) - u(s_{\underline{j}+1}^2; t_M)] \geq 0$$

because  $t_M < \theta_{\underline{j}} < s_{\underline{j}+1}^2$ .

Suppose  $\underline{j} = 1$ . Then, for every  $t > t_M$ ,

$$\alpha(t) = \beta(t) = \min \left( \arg \max \left\{ u(\theta_i; t) | i \in \arg \max \{u(s_j^2; t) | j = 2, \dots, \ell\} \right\} \right). \quad (40)$$

This implies that  $T_2(\hat{v}') = \{t \in T | t \in [t_M, \bar{y}_2]\}$ , and so  $b_2(\tilde{v}) = b_2(\hat{v}')$ . Also, for every

$i = 3, \dots, \ell$ ,  $b_i(\hat{v}') = b_i(\tilde{v})$ . Thus,  $\hat{\Delta}(t_M|0) + \tilde{\Delta}(t_M|0) = 0$ .

Suppose  $\theta_1 < q < \theta_2$ . Then, for every  $i = 2, \dots, \ell$ ,  $s_i^1 = q$ , which implies that, for every  $t > t_M$ , (40) is true. Then,  $b_2(\tilde{v}) = b_2(\hat{v}) + 1$ , and, for every  $i = 3, \dots, \ell$ ,  $b_i(\hat{v}) = b_i(\tilde{v})$ . Thus,  $\Delta(t_M|0) + \tilde{\Delta}(t_M|0) = 0$ .

Lastly, suppose  $q > \theta_2$ . Then, for every  $i = 2, \dots, \ell$ ,  $s_i^1 = s_i^2$ . Then, again, for every  $t > t_M$ , (40) is true, implying  $\hat{\Delta}(t_M|0) + \tilde{\Delta}(t_M|0) = 0$ .

CASE 2: Assume  $m = \ell - 1$ . A symmetric argument can prove the statement for this case.

CASE 3: Assume  $1 < m < \ell - 1$ .

Party  $m$  is the median party in  $\hat{v}$  and  $\tilde{v}'$ , and party  $m + 1$  is the median party in  $\hat{v}'$  and  $\tilde{v}$ . Then,

$$\hat{\Delta}(t_M|0) = \frac{1}{n} \left( \sum_{i \in L} b_i(\hat{v}) [u(s_i^m; t_M) - u(s_i^{m+1}; t_M)] + u(s_m^{m+1}; t_M) - u(s_{m+1}^{m+1}; t_M) \right) \quad (41)$$

and

$$\tilde{\Delta}(t_M|0) = \frac{1}{n} \left( \sum_{i \in L} b_i(\tilde{v}) [u(s_i^{m+1}; t_M) - u(s_i^m; t_M)] + u(s_{m+1}^m; t_M) - u(s_m^m; t_M) \right). \quad (42)$$

By adding (41) and (42), we obtain

$$\begin{aligned} \hat{\Delta}(t_M|0) + \tilde{\Delta}(t_M|0) &= \frac{1}{n} \left( \sum_{i \in L} [b_i(\hat{v}) - b_i(\tilde{v})] [u(s_i^m; t_M) - u(s_i^{m+1}; t_M)] \right. \\ &\quad \left. + u(s_m^{m+1}; t_M) - u(s_{m+1}^{m+1}; t_M) + u(s_{m+1}^m; t_M) - u(s_m^m; t_M) \right). \end{aligned} \quad (43)$$

First, assume  $q < \theta_m$ . Let  $L^- = \{i \in L | \theta_i < 2\theta_m - q\}$ . Note that, for every  $i \in L^-$ ,  $s_i^m = s_i^{m+1}$ . In particular,  $m \in L^-$ . Also, if  $i \notin L^-$ , then  $s_i^m = 2\theta_m - q$ . Then, (43) is reduced



to

$$\begin{aligned} \hat{\Delta}(t_M|0) + \tilde{\Delta}(t_M|0) &= \frac{1}{n} \left( \sum_{i \in L \setminus L^-} [b_i(\hat{v}) - b_i(\tilde{v})][u(2\theta_m - q; t_M) - u(s_i^{m+1}; t_M)] \right. \\ &\quad \left. + u(s_{m+1}^m; t_M) - u(\theta_{m+1}; t_M) \right) \end{aligned} \quad (44)$$

Suppose  $L^- = L$ . Then,  $s_{m+1}^m = \theta_{m+1}$ , so  $\hat{\Delta}(t_M|0) + \tilde{\Delta}(t_M|0) = 0$ . Suppose  $L^- \neq L$ . Let  $\underline{j} = \max L^-$ . Suppose  $\underline{j} \geq m + 1$ . If  $t \leq \frac{\theta_{\underline{j}} + 2\theta_m - q}{2}$ , then  $\hat{v}_t \in L^-$  and  $\tilde{v}_t \in L^-$ . If  $t > \frac{\theta_{\underline{j}} + s_{\underline{j}+1}^{m+1}}{2}$ , then

$$\alpha(t) = \beta(t) = \min \left( \arg \max \left\{ u(\theta_i; t) \mid i \in \arg \max \{ u(s_j^{m+1}; t) \mid j = \underline{j} + 1, \dots, \ell \} \right\} \right),$$

so  $\hat{v}_t = \tilde{v}_t$ . Let  $\tilde{T} = \{t \in T \mid \frac{\theta_{\underline{j}} + 2\theta_m - q}{2} < t \leq \frac{\theta_{\underline{j}} + s_{\underline{j}+1}^{m+1}}{2}\}$ . For every  $t \in \tilde{T}$ ,  $\hat{v}_t = \underline{j} + 1$  and  $\tilde{v}_t = \underline{j}$ . This implies that for every  $i > \underline{j} + 1$ ,  $b_i(\hat{v}) = b_i(\tilde{v})$ , and  $b_{\underline{j}+1}(\hat{v}) - b_{\underline{j}+1}(\tilde{v}) = |\tilde{T}|$ . Then, from (44), we have

$$\hat{\Delta}(t_M|0) + \tilde{\Delta}(t_M|0) = \frac{|\tilde{T}|}{n} \left[ u(2\theta_m - q; t_M) - u(s_{\underline{j}+1}^{m+1}; t_M) \right] \geq 0,$$

because  $t_M < 2\theta_m - q < s_{\underline{j}+1}^{m+1}$ . Now suppose  $\underline{j} = m$ . For every  $t \leq \frac{\theta_{m+1} + s_{m+2}^{m+1}}{2}$ ,  $\alpha(t) \leq m + 1$  and  $\beta(t) \leq m + 1$ . For every  $t > \frac{\theta_{m+1} + s_{m+2}^{m+1}}{2}$ ,

$$\alpha(t) = \beta(t) = \min \left( \arg \max \left\{ u(\theta_i; t) \mid i \in \arg \max \{ u(s_j^{m+1}; t) \mid j = m + 2, \dots, \ell \} \right\} \right),$$

so  $\hat{v}_t = \tilde{v}_t$ . This implies that for every  $t > m + 1$ ,  $T_i(\hat{v}) = T_i(\tilde{v})$ , so  $b_i(\hat{v}) = b_i(\tilde{v})$ . Also, from the strategies,  $T_{m+1}(\hat{v}) = \{t \in T \mid \alpha(t) = m + 1\} = \{t \in T \mid t_{M+1} \leq t \leq \bar{y}_{m+1}\}$ , and  $T_{m+1}(\tilde{v}) = \{t \in T \mid t_M \leq t \leq \bar{y}_{m+1}\}$ , implying  $b_{m+1}(\hat{v}) - b_{m+1}(\tilde{v}) = -1$ . Then, from (44), we conclude that  $\hat{\Delta}(t_M|0) + \tilde{\Delta}(t_M|0) = 0$ .

Second, assume  $\theta_m < q < \theta_{m+1}$ . Then, for every  $i = 1, \dots, m$ ,  $s_i^{m+1} = q$ , and, for every

$i = m + 1, \dots, \ell$ ,  $s_i^m = q$ . Then from (43) we have

$$\begin{aligned} \hat{\Delta}(t_M|0) + \tilde{\Delta}(t_M|0) &= \frac{1}{n} \left( \sum_{i=1}^m [b_i(\hat{v}) - b_i(\tilde{v})][u_i(s_i^m; t_M) - u_i(q; t_M)] \right. \\ &\quad + \sum_{i=m+1}^{\ell} [b_i(\hat{v}) - b_i(\tilde{v})][u_i(q; t_M) - u_i(s_i^{m+1}; t_M)] \\ &\quad \left. + 2u(q; t_M) - u(\theta_m; t_M) - u(\theta_{m+1}; t_M) \right). \end{aligned} \quad (45)$$

For every  $t < \underline{y}_m$ ,

$$\alpha(t) = \beta(t) = \min \left( \arg \max \left\{ u(\theta_i; t) \mid i \in \arg \max \{ u(s_j^m; t) \mid j = 1, \dots, m-1 \} \right\} \right).$$

For every  $t > \bar{y}_{m+1}$ ,

$$\alpha(t) = \beta(t) = \min \left( \arg \max \left\{ u(\theta_i; t) \mid i \in \arg \max \{ u(s_j^{m+1}; t) \mid j = m+2, \dots, \ell \} \right\} \right).$$

Also, for every  $t \in [\underline{y}_m, \bar{y}_{m+1}]$ ,  $\{\hat{v}_t, \tilde{v}_t\} = \{m, m+1\}$ . Therefore, for every  $i \in L \setminus \{m, m+1\}$ ,  $T_i(\hat{v}) = T_i(\tilde{v})$ , implying  $b_i(\hat{v}) = b_i(\tilde{v})$ . Note that  $t_{m+1} > \bar{y}_m$ . So,  $T_m(\hat{v}) = \{t \in T \mid t \in [\underline{y}_m, t_M]\}$  and  $T_m(\tilde{v}) = \{t \in T \mid t \in [\underline{y}_m, t_{M-1}]\}$ . This implies  $b_m(\hat{v}) - b_m(\tilde{v}) = 1$ . Also,  $T_{m+1}(\hat{v}) = \{t \in T \mid t \in [t_{m+1}, \bar{y}_{m+1}]\}$  and  $T_{m+1}(\tilde{v}) = \{t \in T \mid t \in [t_M, \bar{y}_{m+1}]\}$ , implying  $b_{m+1}(\hat{v}) - b_{m+1}(\tilde{v}) = -1$ . Then, from (45), we conclude that  $\hat{\Delta}(t_M|0) + \tilde{\Delta}(t_M|0) = 0$ .

Lastly, assume  $q > \theta_{m+1}$ . Let  $L^+ = \{i \in L \mid \theta_i > 2\theta_{m+1} - q\}$ . Then, for every  $i \in L^+$ ,  $s_i^m = s_i^{m+1}$ , and in particular  $m+1 \in L^+$ . For every  $i \notin L^+$ ,  $s_i^{m+1} = 2\theta_{m+1} - q$ . Then, we have

$$\begin{aligned} \hat{\Delta}(t_M|0) + \tilde{\Delta}(t_M|0) &= \frac{1}{n} \left( \sum_{i \in L \setminus L^+} [b_i(\hat{v}) - b_i(\tilde{v})][u(s_i^m; t_M) - u(2\theta_{m+1} - q; t_M)] \right. \\ &\quad \left. + u(s_m^{m+1}; t_M) - u(\theta_m; t_M) \right). \end{aligned} \quad (46)$$

First, if  $L^+ = L$ , then clearly  $\hat{\Delta}(t_M|0) + \tilde{\Delta}(t_M|0) = 0$ . Suppose  $L^+ \neq L$ . Let  $\underline{j} = \min L^+$ . Suppose  $\underline{j} \leq m$ . Then, if  $t \leq \frac{\theta_{\underline{j}} + s_{\underline{j}-1}^m}{2}$ , then  $\alpha(t) = \beta(t) \leq \underline{j} - 1$ . Let  $\tilde{T} = \{t \in T | \frac{\theta_{\underline{j}} + s_{\underline{j}-1}^m}{2} < t \leq \frac{\theta_{\underline{j}} + 2\theta_{m+1} - q}{2}\}$ . If  $t \in \tilde{T}$ , then  $\hat{v}_t = \alpha(t) = \underline{j}$  and  $\tilde{v}_t = \beta(t) = \underline{j} - 1$ . Then, for every  $i < \underline{j} - 1$ ,  $b_i(\hat{v}) - b_i(\tilde{v}) = 0$ , and  $b_{\underline{j}-1}(\hat{v}) - b_{\underline{j}-1}(\tilde{v}) = -|\tilde{T}|$ . Then, from (46), we have

$$\hat{\Delta}(t_M|0) + \tilde{\Delta}(t_M|0) = -\frac{|\tilde{T}|}{n} \left[ u(s_{\underline{j}-1}^m; t_M) - u(2\theta_{m+1} - q; t_M) \right] \geq 0$$

because  $s_{\underline{j}-1}^m > 2\theta_{m+1} - q > \theta_m > t_M$ . Suppose  $\underline{j} = m + 1$ . If  $t < \underline{y}_m$ , then  $\alpha(t) = \beta(t)$ . Thus, for every  $i < m$ ,  $b_i(\hat{v}) - b_i(\tilde{v}) = 0$ . From the strategies,  $T_m(\hat{v}) = \{t \in T | t \in [\underline{y}_m, t_M]\}$  and  $T_m(\tilde{v}) = \{t \in T | t \in [\underline{y}_m, t_{M+1}]\}$ . So,  $b_m(\hat{v}) - b_m(\tilde{v}) = 1$ . Then, clearly from (46)  $\hat{\Delta}(t_M|0) + \tilde{\Delta}(t_M|0) = 0$ .  $\blacksquare$

Lemma 6 implies that either  $\hat{\Delta}(t_M|0) > 0$ , or  $\tilde{\Delta}(t_M|0) > 0$ , or  $\hat{\Delta}(t_M|0) = \tilde{\Delta}(t_M|0) = 0$ . Then, by Lemma 5, either  $\hat{v}$  or  $\tilde{v}$  is a robust equilibrium of  $G(T, \theta, q, 0)$ .

Suppose  $\hat{v} \in V(T, \theta, q)$ . By construction, for every  $t \notin [\underline{y}_m, t_M]$ ,  $\hat{v}_t = \alpha(t)$ , so  $\hat{v}_t$  is strategically sincere in  $\hat{v}$ . For any  $t \in [\underline{y}_m, t_M]$ ,  $\hat{v}_t = m$ , and  $m \notin \arg \max\{u(s_i^m; t) | i \in L\}$  if and only if  $t > \bar{y}_m$ . Therefore,  $T \setminus T^*(\hat{v}) = \{t \in T | \bar{y}_m < t \leq t_M\}$ . Suppose  $\tilde{v}$  is a robust equilibrium. By construction, for every  $t \notin [t_M, \bar{y}_{m+1}]$ ,  $\tilde{v}_t = \beta(t)$ , so  $\tilde{v}_t$  is strategically sincere in  $\tilde{v}$ . For any  $t \in [t_M, \bar{y}_{m+1}]$ ,  $\tilde{v}_t = m + 1$ , and  $m + 1 \notin \arg \max\{u(s_i^{m+1}; t) | i \in L\}$  if and only if  $t < \underline{y}_{m+1}$ . Therefore,  $T \setminus T^*(\tilde{v}) = \{t \in T | t_M \leq t < \underline{y}_{m+1}\}$ .  $\blacksquare$

## Proof of Proposition 5

Assume that  $v$  and  $v'$  are strategically sincere robust equilibria of  $G(T, \theta, q)$ . By Proposition 6,  $k(v) = k(v') = m$ . Then, since  $v$  and  $v'$  are strategically sincere and A2 holds,  $b_m(v) = b_m(v') = |X_m|$ . If  $|X_m| \geq M$ , then  $\lambda^v = \lambda^{v'}$  as both of them are the degenerate lottery on

$\theta_m$ . Suppose  $|X_m| < M$ . Then, for each  $x \in \mathbb{R}$ ,

$$\lambda^v(x) = \frac{1}{n} \sum_{\{i \in L | s_i^m = x\}} b_i(v) \quad \text{and} \quad \lambda^{v'}(x) = \frac{1}{n} \sum_{\{i \in L | s_i^m = x\}} b_i(v').$$

Since A2 holds, and  $v$  and  $v'$  are strategically sincere, for every  $x$  with  $\{i \in L | s_i^m = x\} \neq \emptyset$ ,

$$\frac{1}{n} \sum_{\{i \in L | s_i^m = x\}} b_i(v) = \frac{1}{n} \sum_{\{i \in L | s_i^m = x\}} b_i(v') = \left| \left\{ t \in T \mid u(x; t) = \max\{u(s_j^m; t) \mid j \in L\} \right\} \right|,$$

which completes the proof. ■

## Proof of Lemma 2

Let  $v \in V(T, \theta, q)$  and let  $t \in T \setminus T^*(v)$ . Let  $k = k(v)$  and  $i = v_t$ . Then,  $i \notin \arg \max\{u(s_h^k; t) \mid h \in L\}$ . Suppose  $i \neq k$ . Since  $L$  is finite,  $\arg \max\{u(s_h^k; t) \mid h \in L\} \neq \emptyset$ . Let  $j \in \arg \max\{u(s_h^k; t) \mid h \in L\}$  and let  $v' = (j, v_{-t})$ . First, suppose  $k(v') = k$ . Then,  $p_h(v') = s_h^k$  for every  $h \in L$ . If  $k$  is not the majority party in both  $v$  and  $v'$ , then, for every  $\epsilon \in [0, 1]$

$$U(v; t|\epsilon) - U(v'; t|\epsilon) = \frac{1}{n} [u(s_i^k; t) - u(s_j^k; t)] < 0,$$

contradicting that  $v$  is a robust equilibrium. If  $k$  is the majority party in both  $v$  and  $v'$ , then, for every  $\epsilon \in (0, 1]$ ,

$$U(v; t|\epsilon) - U(v'; t|\epsilon) = \frac{\epsilon}{n} [u(s_i^k; t) - u(s_j^k; t)] < 0,$$

a contradiction. If  $k$  is not the majority party in  $v$ , but it is in  $v'$ , then it must be the case that  $b_k(v) = M - 1$  and  $j = k$ . Then, for every  $\epsilon \in [0, 1]$ ,

$$U(v; t|\epsilon) - U(v'; t|\epsilon) = (1 - \epsilon) \left[ \sum_{h \in L} \frac{b_h(v)}{n} u(s_h^k; t) - u(s_j^k; t) \right] + \frac{\epsilon}{n} [u(s_i^k; t) - u(s_j^k; t)] < 0,$$

a contradiction.

Secondly, suppose  $k(v') \neq k$ . Suppose  $i < k$ . Then, it must be the case that  $\sum_{h=1}^k b_h(v) = M$  and  $j > k$ . Since  $j > k$ ,  $t \geq \bar{y}_k$ . Then, since  $s_i^k < s_k^k = \theta_k < t$ , we have  $u(s_i^k; t) < u(s_k^k; t)$ . I also claim that  $b_k(v) < M - 1$ . To see this, suppose  $b_k(v) = M - 1$ . Note that  $b_i(v) + b_k(v) = M$  and  $t \neq t_1$ . This implies  $v_{t_1} \geq k > 1$ . But since  $t_1 < \theta_1$ ,  $u(s_1^k; t_1) > u(s_{v_{t_1}}^k; t_1)$ . Also, party  $k$  would remain as the median party even after voter  $t_1$ 's deviation by voting for party 1. Then, for every  $\epsilon \in [0, 1)$ ,  $U(v; t_1|\epsilon) < U(1, v_{-t_1}; t_1|\epsilon)$ , a contradiction that implies that the claim is true. Then, for every  $\epsilon \in [0, 1)$ ,

$$U(v; t|\epsilon) - U(k, v_{-t}; t|\epsilon) = \frac{1}{n}[u(s_i^k; t) - u(s_k^k; t)] < 0,$$

a contradiction. A symmetric argument will lead to a contradiction when  $i > k$ . ■

## Proof of Proposition 6

Let  $v$  be a strategically sincere robust equilibrium. Let  $k = k(v)$ . Since  $v$  is strategically sincere, for every  $t < \underline{y}_k$ ,  $v_t < k$ ; and for every  $t > \bar{y}_k$ ,  $v_t > k$ . Then, for  $k$  to be decisive, it must be that  $t_M \in X_k$ . Since  $t_M \in [\theta_m, \frac{\theta_m + \theta_{m+1}}{2}]$ , either  $k = m$  or  $k = m + 1$ . Suppose  $k = m + 1$ . Then,  $t_M = \frac{\theta_m + \theta_{m+1}}{2}$  and  $\underline{x}_{m+1}(q) = \theta_m$ . Since  $s_m^{m+1} = \theta_m$  and  $s_{m+1}^{m+1} = \theta_{m+1}$ , we have  $\max\{u(s_i^{m+1}; t_M) | i \in L\} = u(s_m^{m+1}; t_M) = u(s_{m+1}^{m+1}; t)$ , contradicting A2. Thus,  $k(v) = m$ . ■

## Proof of Proposition 7

Let  $v \in V(T, \theta, q)$ . Suppose  $v$  is strategically sincere and satisfies C1. Suppose  $v_t$  is strategic. Let  $j = v_t$  and  $k = k(v)$ . Suppose  $t \in (\underline{y}_k, \bar{y}_k)$ , then  $j = k$  since  $v$  is strategically sincere. By definition,  $\frac{\theta_{k-1} + \theta_k}{2} \leq \underline{y}_k$  and  $\bar{y}_k \leq \frac{\theta_k + \theta_{k+1}}{2}$ . Then,  $\arg \max\{u(\theta_h; t) | h \in L\} = \{k\}$ , contradicting that  $v_t$  is strategic. Thus, either  $t \leq \underline{y}_k$  or  $t \geq \bar{y}_k$ .

Suppose  $t \leq \underline{y}_k$ . Since  $v$  is strategically sincere,  $j \leq k - 1$ . I claim that  $p_j(v) = \underline{x}_k(q)$ . Suppose not. Then  $p_j(v) = \theta_j > \underline{x}_k(q)$ . Suppose  $t \geq \theta_j$ . Since  $j < k$ ,  $p_{j+1}(v) = \theta_{j+1}$ . Since  $v$  is strategically sincere,  $t \in [\theta_j, \frac{\theta_j + \theta_{j+1}}{2}]$ , implying  $v_t$  is sincere, a contradiction. Suppose  $t < \theta_j$ . If  $j = 1$ , clearly  $\arg \max\{u(\theta_h; t) | h \in L\} = \{1\}$ . So,  $v_t$  is sincere, a contradiction. So,  $j > 1$ . Since  $v$  is strategically sincere,  $t \in [\frac{p_{j-1}(v) + \theta_j}{2}, \theta_j)$ . But  $p_{j-1}(v) = \max\{\theta_{j-1}, \underline{x}_k(q)\} \geq \theta_{j-1}$ , which implies  $t \in [\frac{\theta_{j-1} + \theta_j}{2}, \theta_j)$ . Thus,  $v_t$  is sincere, a contradiction.

Thus, the claim is true,  $p_j(v) = \underline{x}_k(q)$ , which implies  $\theta_j \leq \underline{x}_k(q)$ . I now claim that  $\theta_{j+1} > \underline{x}_k(q)$ . Suppose not. Then,  $p_j(v) = p_{j+1}(v) = \underline{x}_k(q)$ . By C1,  $t \leq \frac{\theta_j + \theta_{j+1}}{2}$ . If  $j = 1$ , then  $v_t$  is sincere, a contradiction. If  $j \geq 2$ , then  $p_{j-1}(v) = \underline{x}_k(q)$ . Then, C1 implies that  $t \geq \frac{\theta_{j-1} + \theta_j}{2}$ . Thus,  $v_t$  is sincere, a contradiction. Hence, the claim is true.

Since  $v$  is strategically sincere,  $t \leq \frac{\underline{x}_k(q) + \theta_{j+1}}{2}$ . If  $t \leq \frac{\theta_j + \theta_{j+1}}{2}$ , then  $v_t$  is sincere. Thus,  $\frac{\theta_j + \theta_{j+1}}{2} < t \leq \frac{\underline{x}_k(q) + \theta_{j+1}}{2} < \theta_{j+1}$ . Then,  $\arg \max\{u(\theta_h; t) | h \in L\} = \{j + 1\}$ . Thus,  $i(t) = j + 1$ , and we have  $\theta_j < t < \theta_{i(t)} \leq \theta_k$ .

I now prove that  $k \geq m$ . Suppose  $k \leq m - 1$ . Since  $v$  is strategically sincere,  $\bigcup_{h=1}^k T_h(v) \subseteq [t_1, \bar{y}_k]$ . But since  $\theta_m \leq t_M$  and  $k \leq m - 1$ ,  $t_M > \bar{y}_k$ . Then  $\sum_{h=1}^k b_h(v) < M$ , contradicting  $k = k(v)$ . Thus,  $k \geq m$ . Therefore,  $\theta_j < t < \theta_{i(t)} \leq \theta_m$ . A symmetric argument will prove that when  $t \geq \bar{y}_k$ , then  $\theta_m \leq \theta_{i(t)} < t < \theta_j$ . ■