

Online appendix for the paper
*Autonomous Agents Coordination: Action
 Languages meet CLP(\mathcal{FD}) and Linda**
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Appendix A The languages \mathcal{B}^{AAC} , B^{MAP} , and B^{MV}

The language \mathcal{B}^{AAC} , and its implementation, heavily relies on its foundations B^{MAP} and B^{MV} . In this section we briefly compare these three languages to clarify which parts of the solvers of the previous languages can be used for the implementation of \mathcal{B}^{AAC} presented in Subsection 3.6.

Let us focus first on B^{MV} . This is a single agent framework. Therefore, considering a given action theory, all fluents and actions are known to the single agent, and the language does not permit to specify private fluents or actions. Moreover, B^{MV} allows one to specify static causal laws. The syntax of fluent expressions and constraints is exactly the same as in \mathcal{B}^{AAC} . The syntax for executability and action effects is analogous to that of \mathcal{B}^{AAC} . More precisely, in B^{MV} , these laws take the forms:

- **executable**(a, C)
- **causes**(x, C_1, C_2), where C_1 is the constraint that will hold in the next state if the action x is executed in a state where C_2 holds.

These are just syntactical variants of (5) and (6), respectively. The semantics of B^{MV} is given via a transition system analogous to that introduced for \mathcal{B}^{AAC} . In particular, one might note that if a \mathcal{B}^{AAC} action description involves a single agent that knows all

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the fluents (and no communication laws are included), then its semantics coincides with the one of the corresponding B^{MV} program obtained by an immediate syntactical translation. The Prolog interpreter for B^{MV} is proved to be correct and complete (for soundness the absence of static laws is needed, but this is the case of \mathcal{B}^{AAC} , as presented here) with respect to the semantics in (Dovier et al. 2010).

Let us consider now B^{MAP} . It is a multiagent, centralized language, where collective actions, namely actions that require more than one agent for being executed, are allowed. For instance, a law of the form

action x executable by a_1, a_2, \dots, a_n

specifies that agents a_1, a_2, \dots, a_n may execute together the action x . In \mathcal{B}^{AAC} , instead, in the domain of an agent a , an action definition implicitly states that the action is executed by a (hence, this is a particular case of the B^{MAP} law). On the other hand, since the reasoner is centralized, conflicts among effects never occur and all (concomitant) planned actions are always successfully executed. The declaration of fluents in B^{MAP} is analogous to that in \mathcal{B}^{AAC} , whereas B^{MAP} has a different syntax for dynamic laws, since they can refer directly to action-occurrences. A B^{MAP} dynamic law has the form *Prec causes Eff*, where *Prec* and *Eff* are constraints and at least one reference to an action x must explicitly occur in *Prec*. Such references are specified by exploiting action flags of the form `actocc(x)`.

The semantics of B^{MAP} is given via the same notion of transition system used for B^{MV} and for \mathcal{B}^{AAC} . If a multi-agent action description in \mathcal{B}^{AAC} , together with initial state and goal, is such that during the plan, no conflict occurs, then the B^{MAP} action description obtained by a simple (mostly one-to-one) translation, has exactly the same behaviour on the transition system. Let us observe that in this translation, collective actions are not generated.